An Empirical Guide to Scalable Persistent Memory

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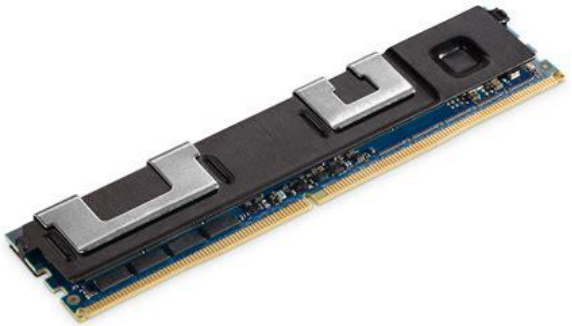
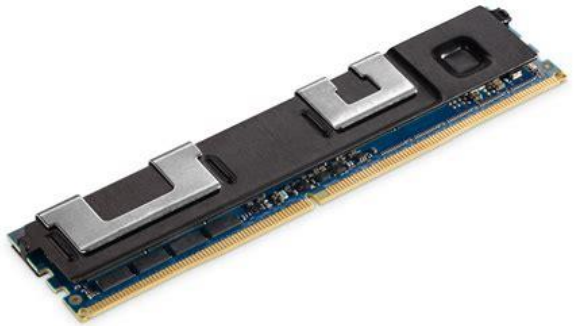
NVSL

Optane DIMMs are Here!



Optane DIMMs are Here!

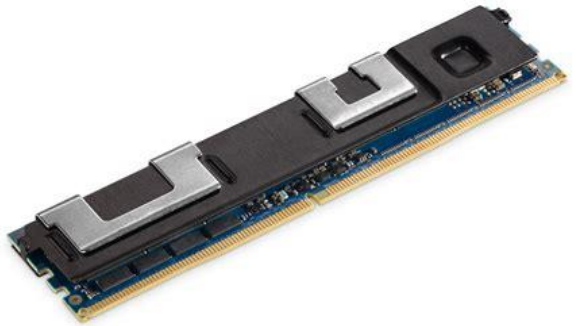
- Not Just Slow Dense DRAM™



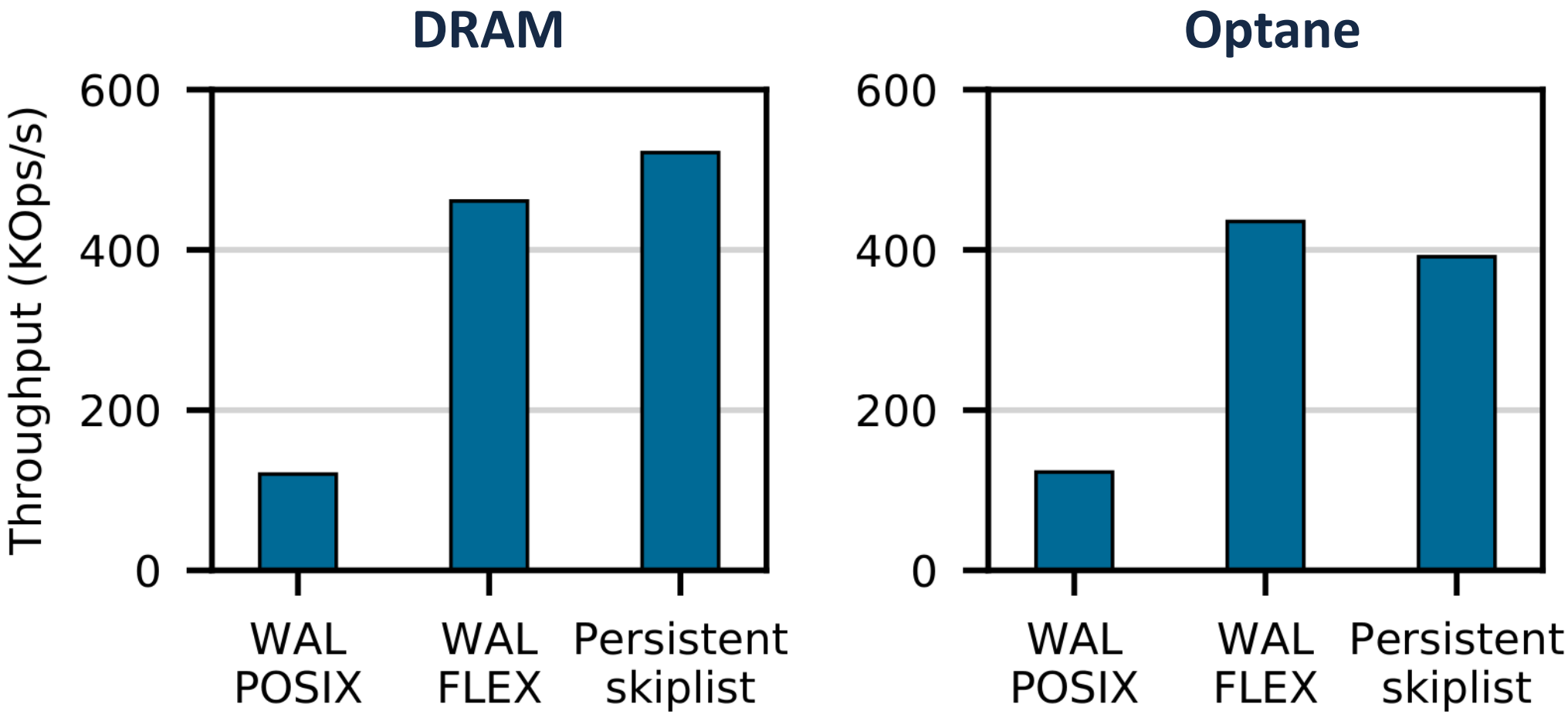
Optane DIMMs are Here!

- Not Just Slow Dense DRAM™

- Slower, denser, media
  - -> More complex architecture
    - -> Second-order performance anomalies
      - -> Fundamentally different

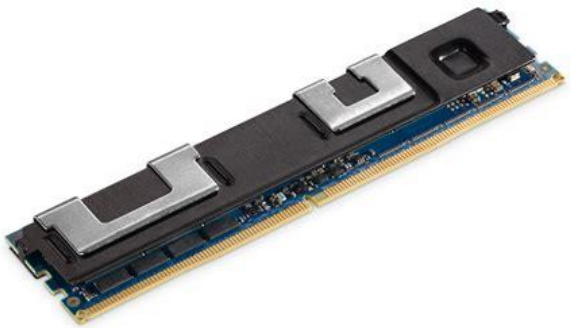


Basics: Emulation is misleading



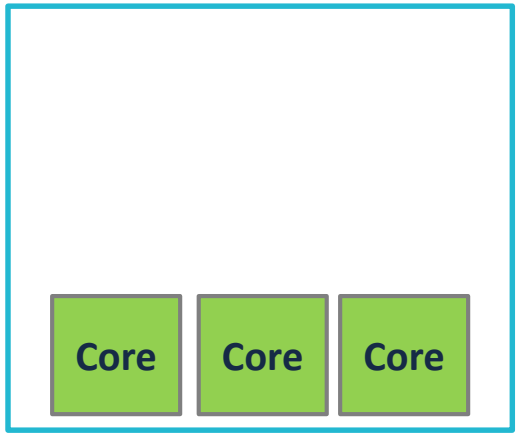
Outline

- Background
- Basics: Optane DIMM Performance
- Lessons: Optane DIMM Best Practices
- Conclusion



Background

Background: Optane in the Machine



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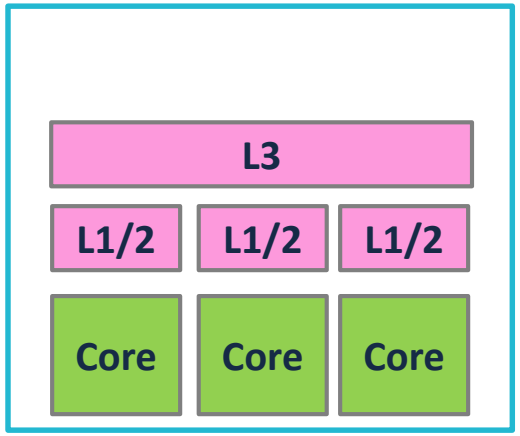

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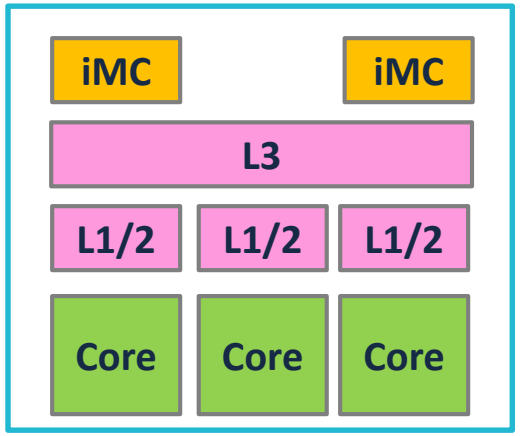

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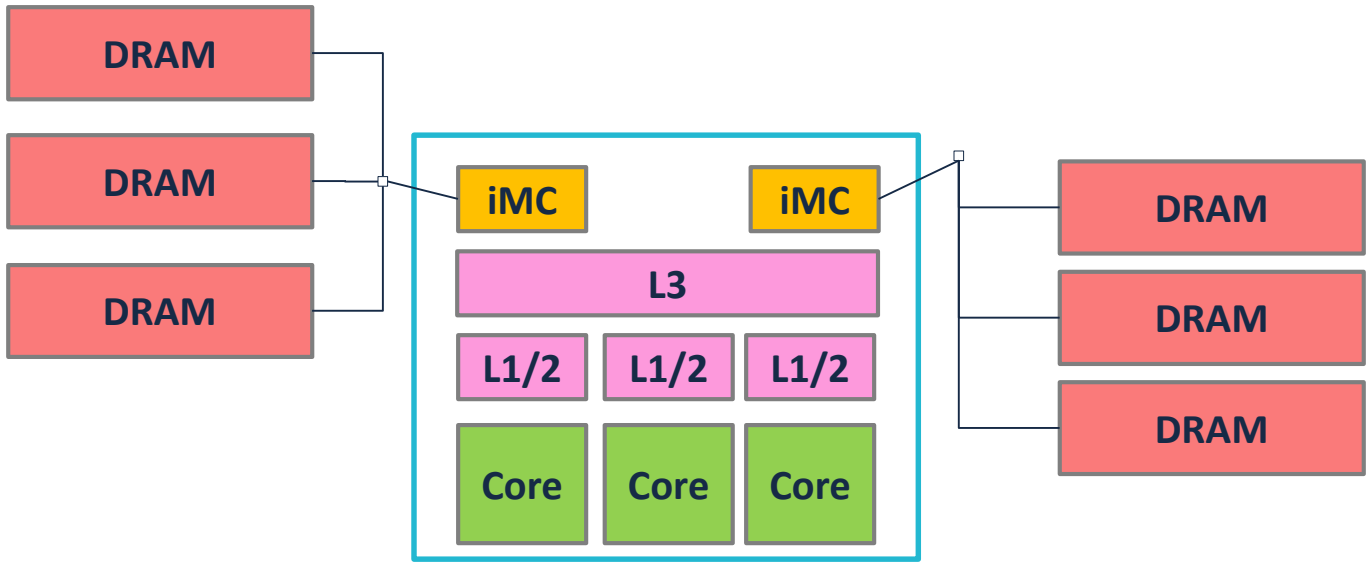

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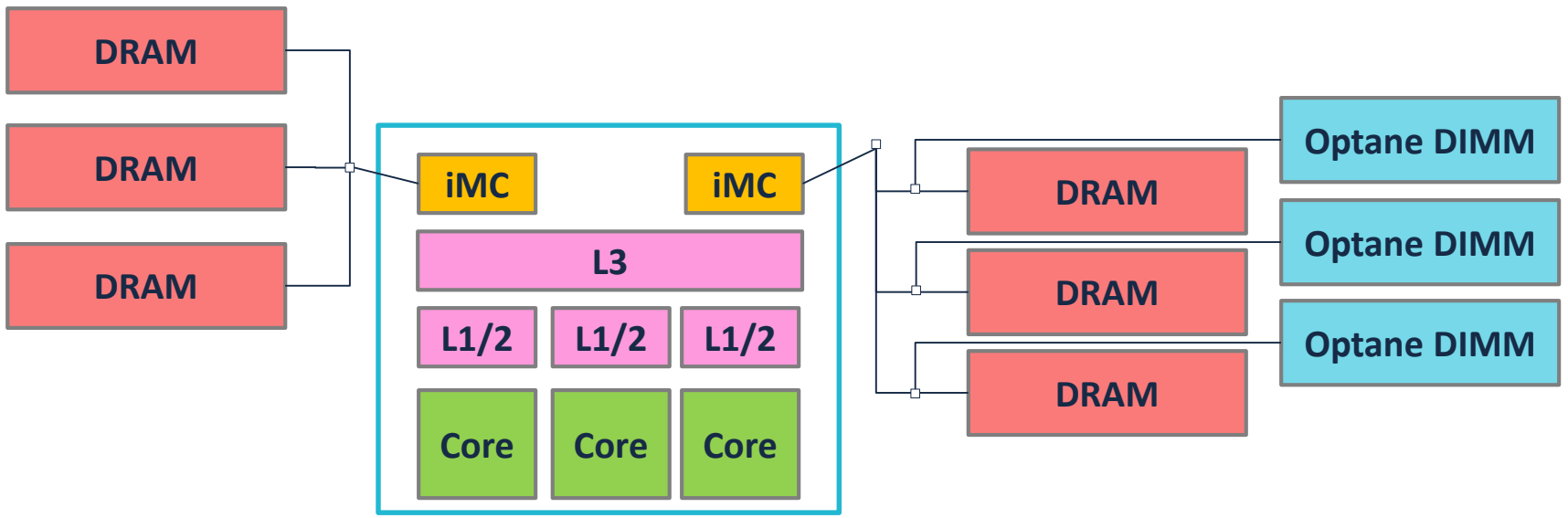

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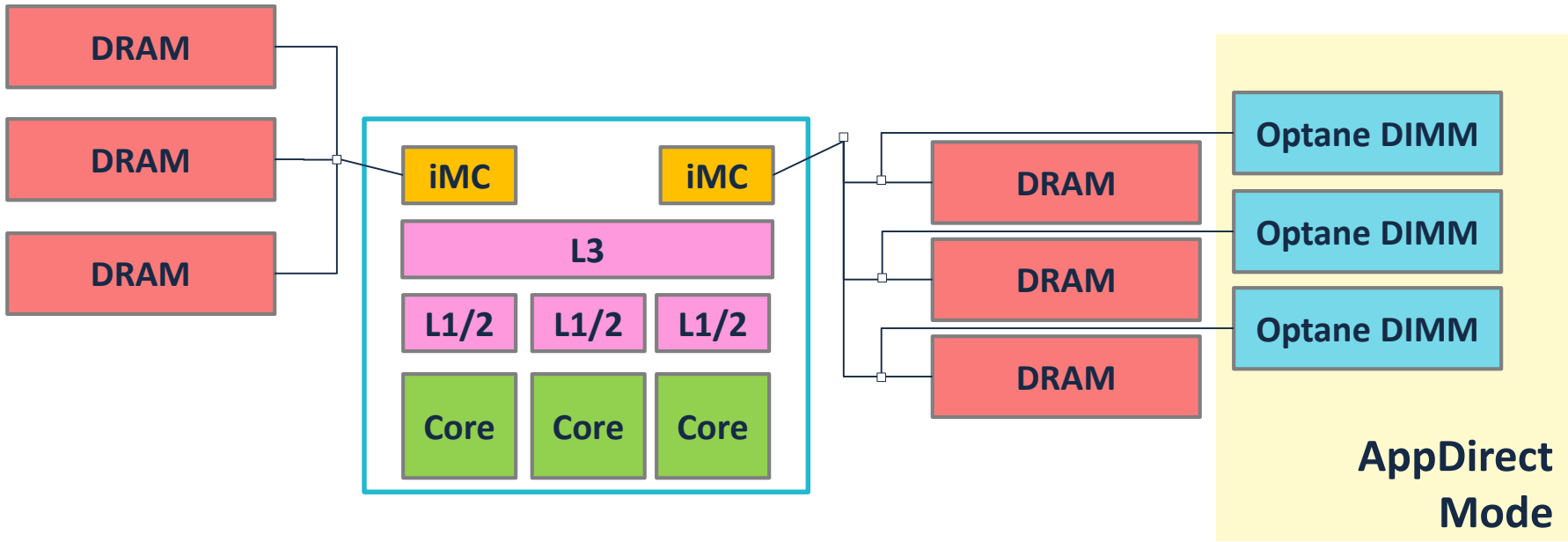

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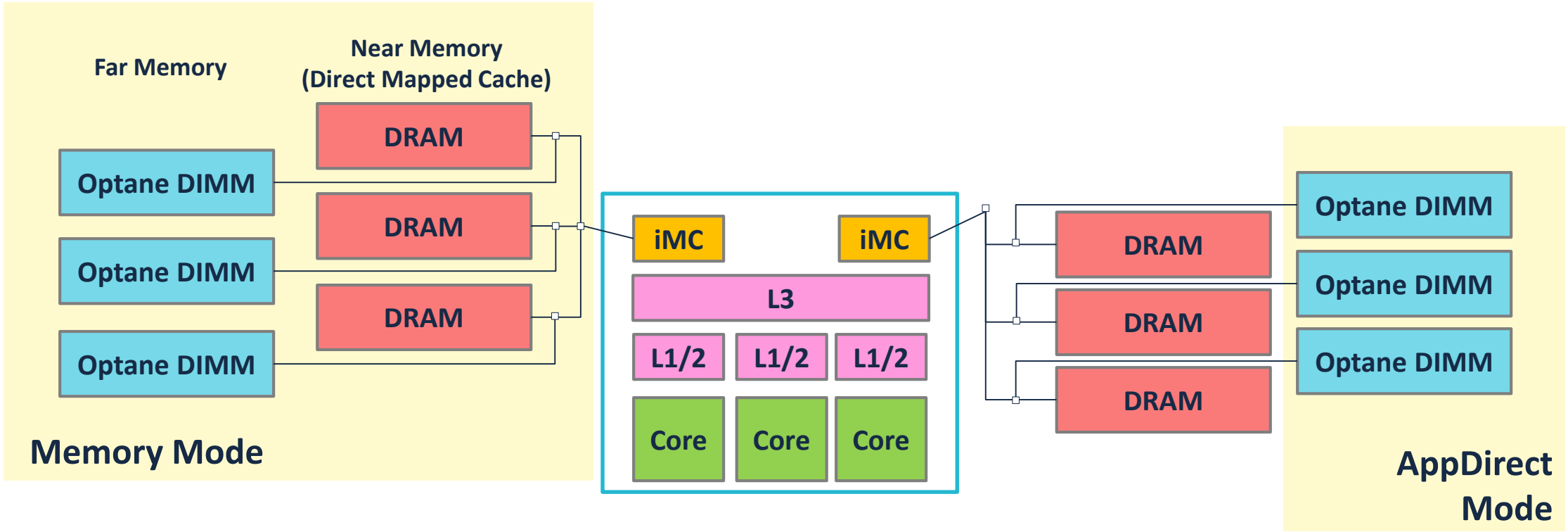

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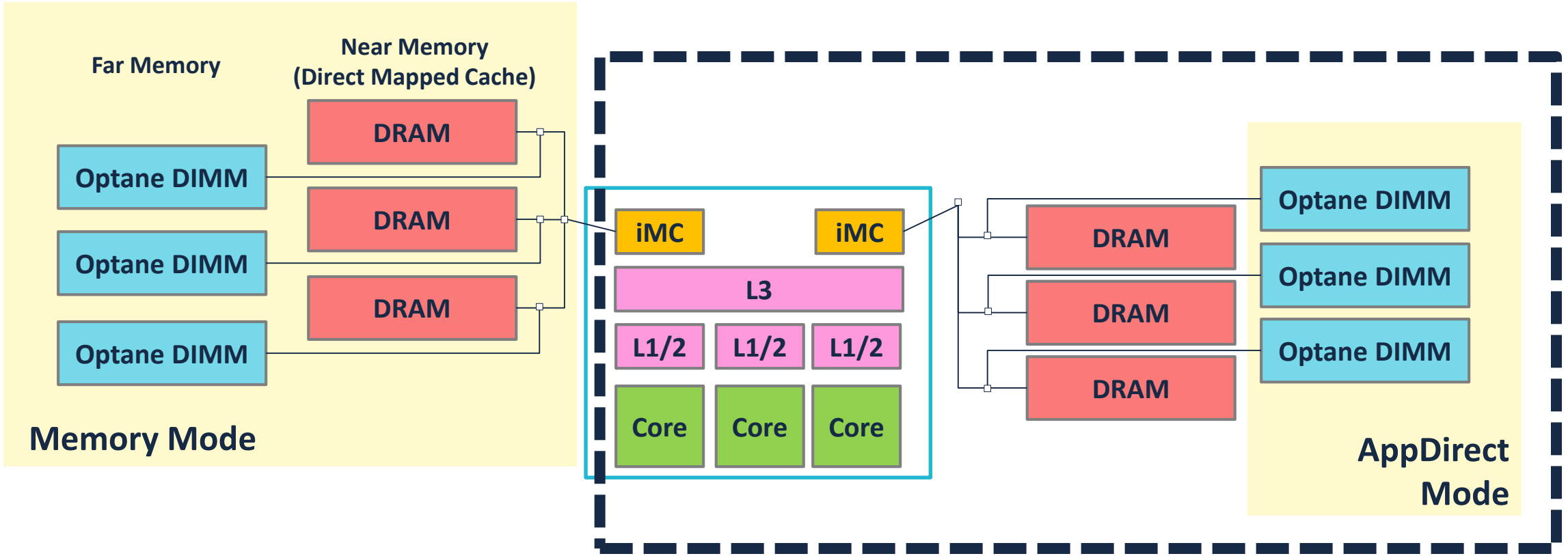
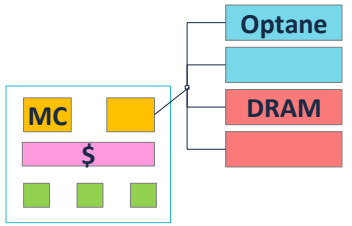


Background: Optane Internals

Optane DIMM



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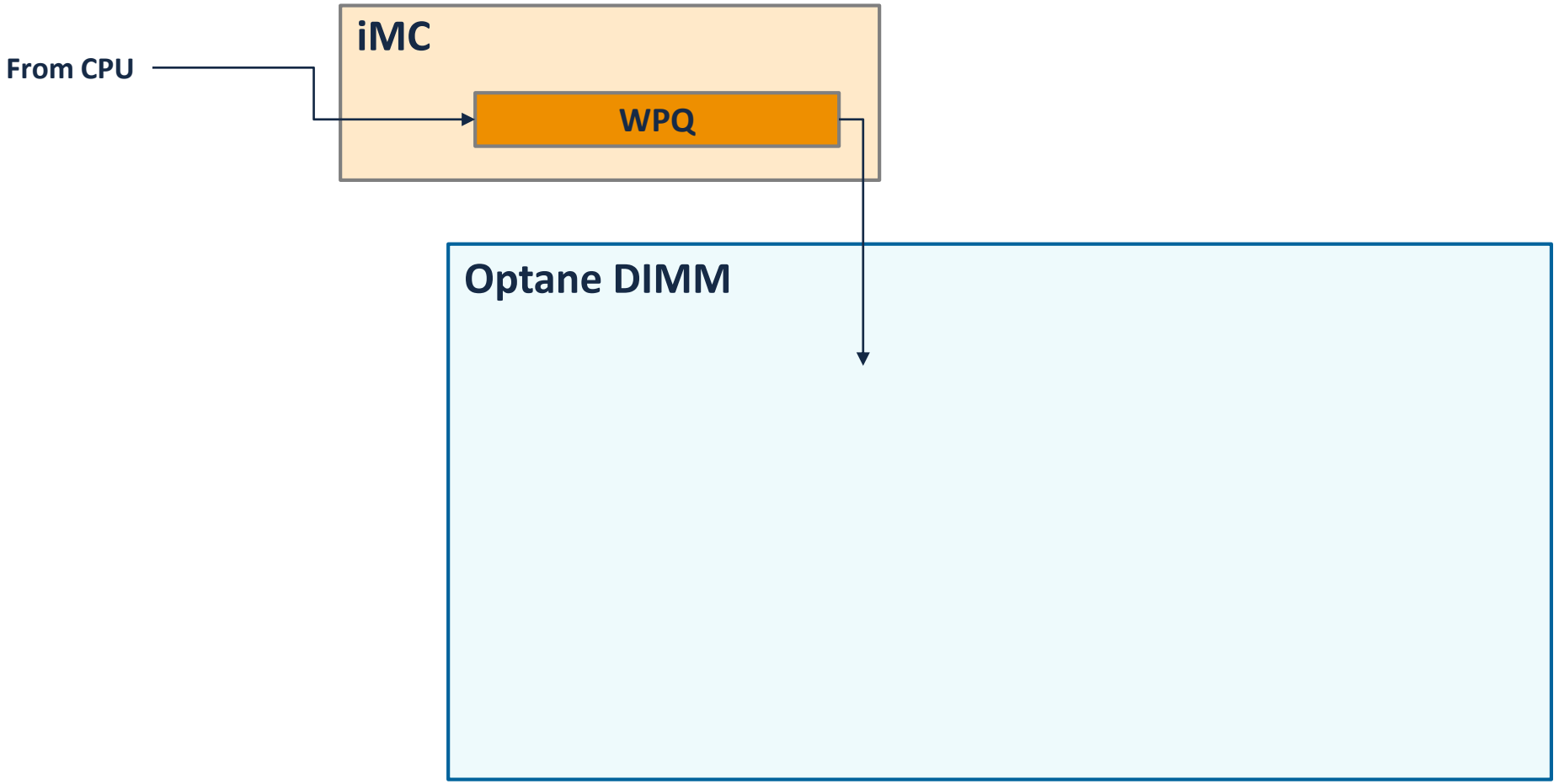

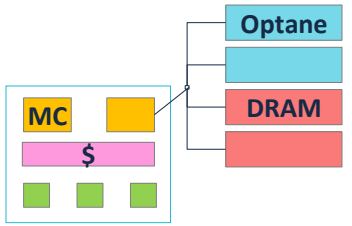

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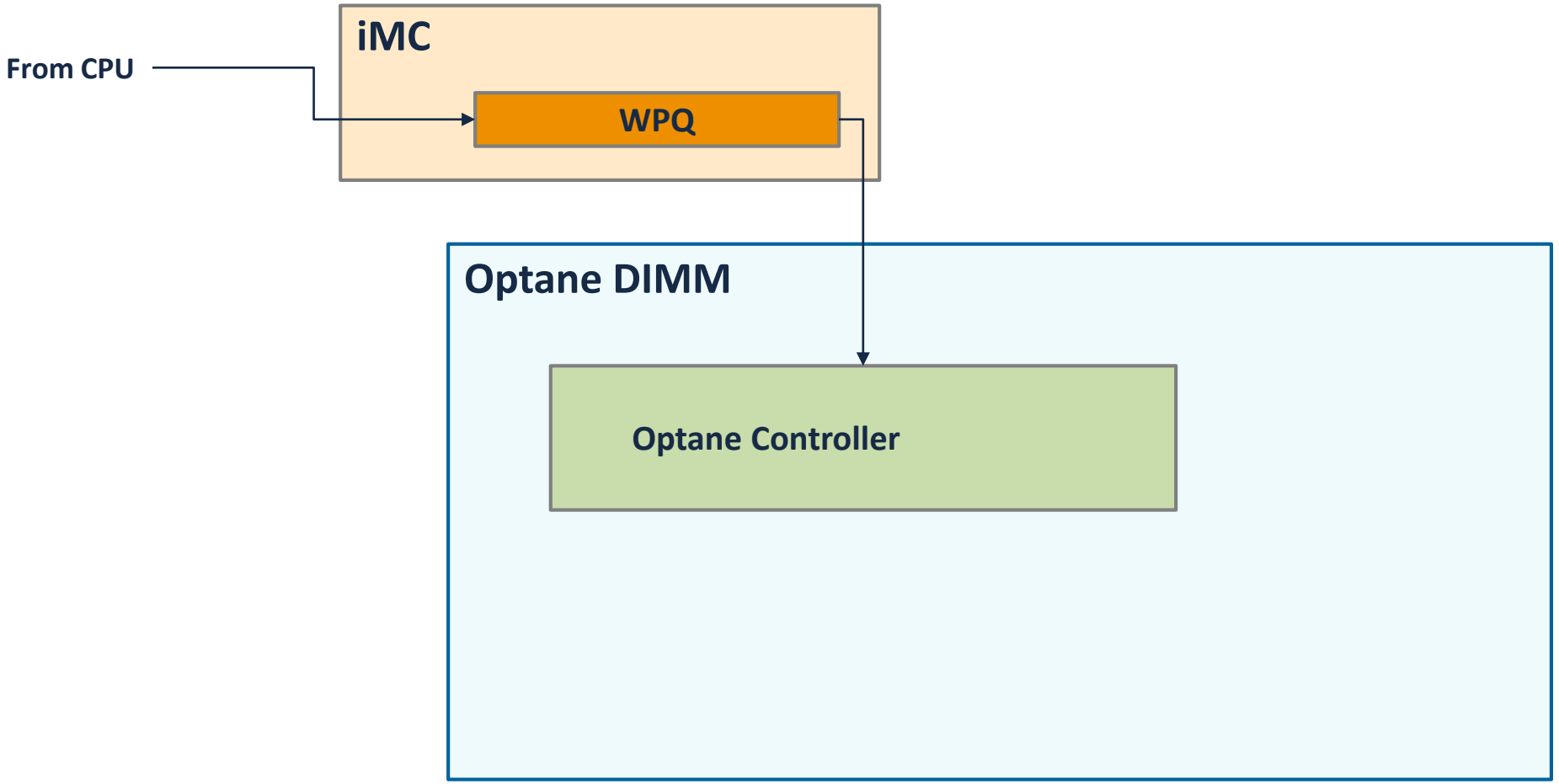

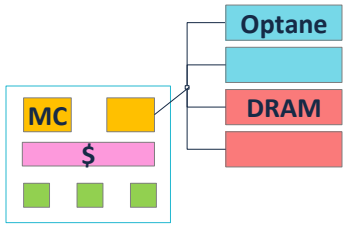

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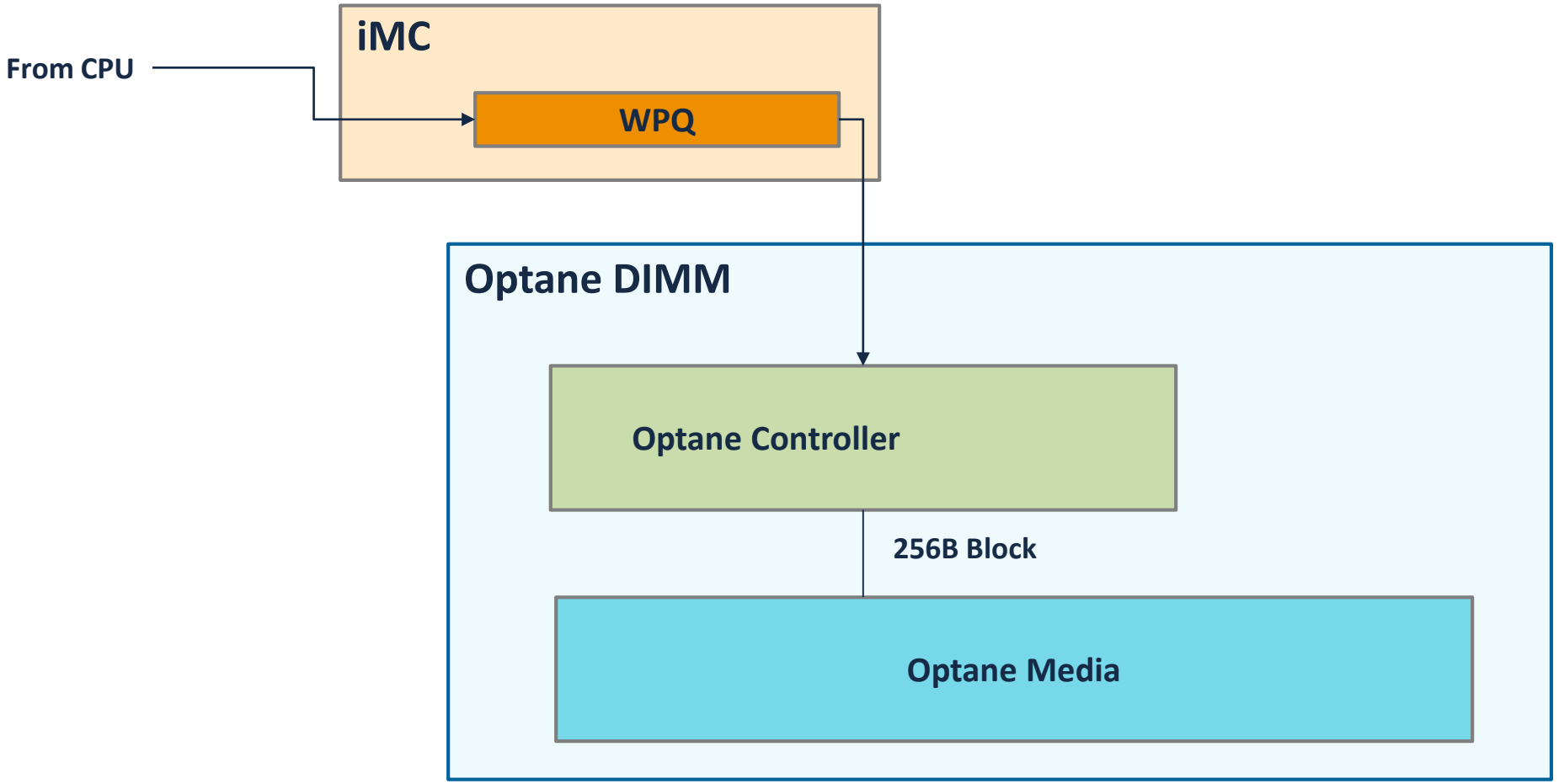

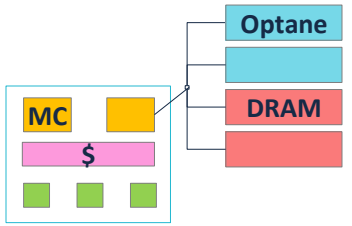

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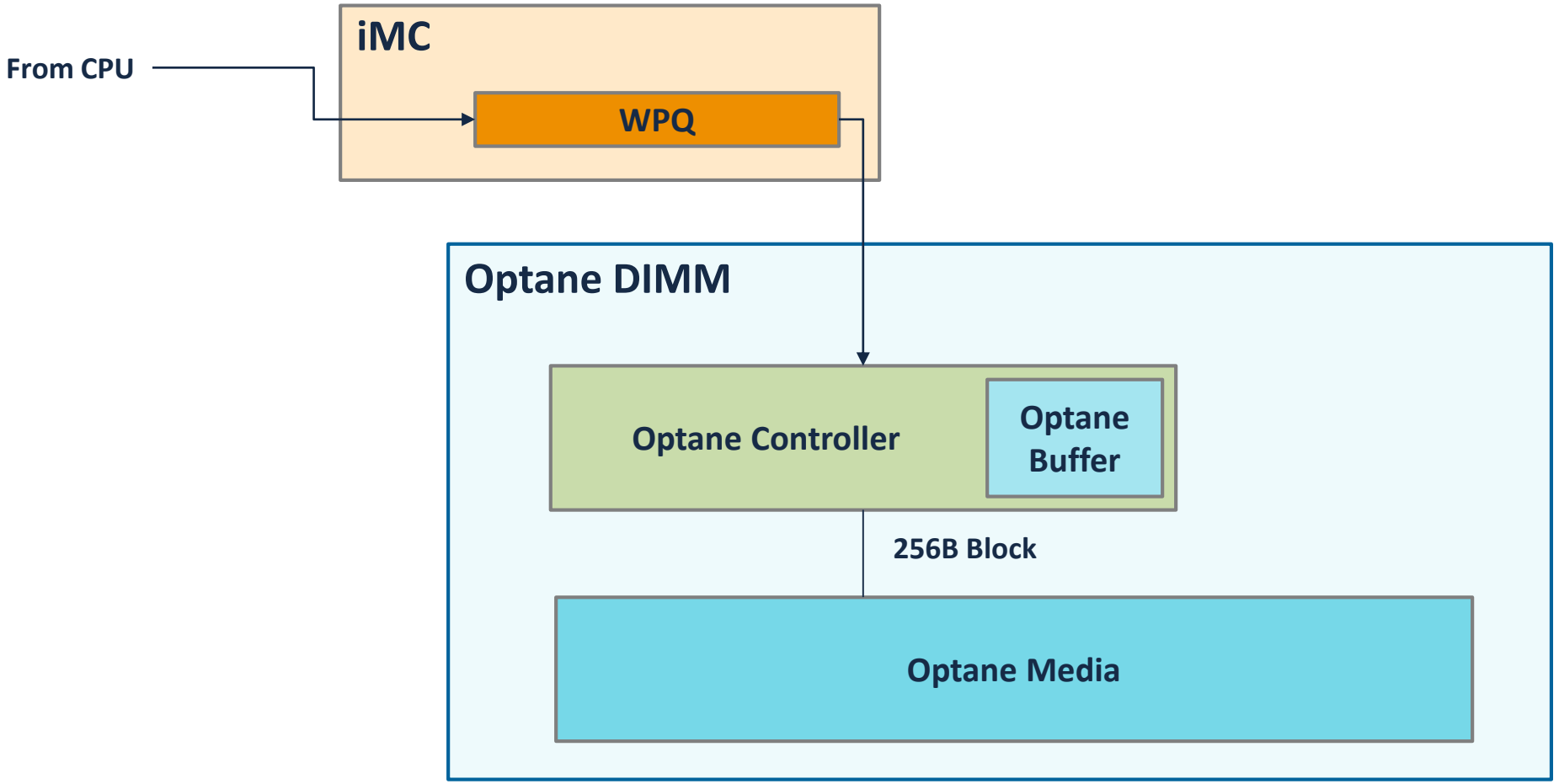

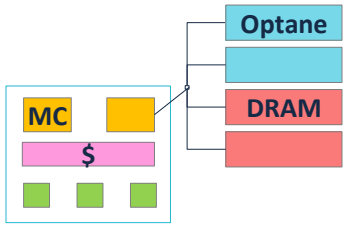

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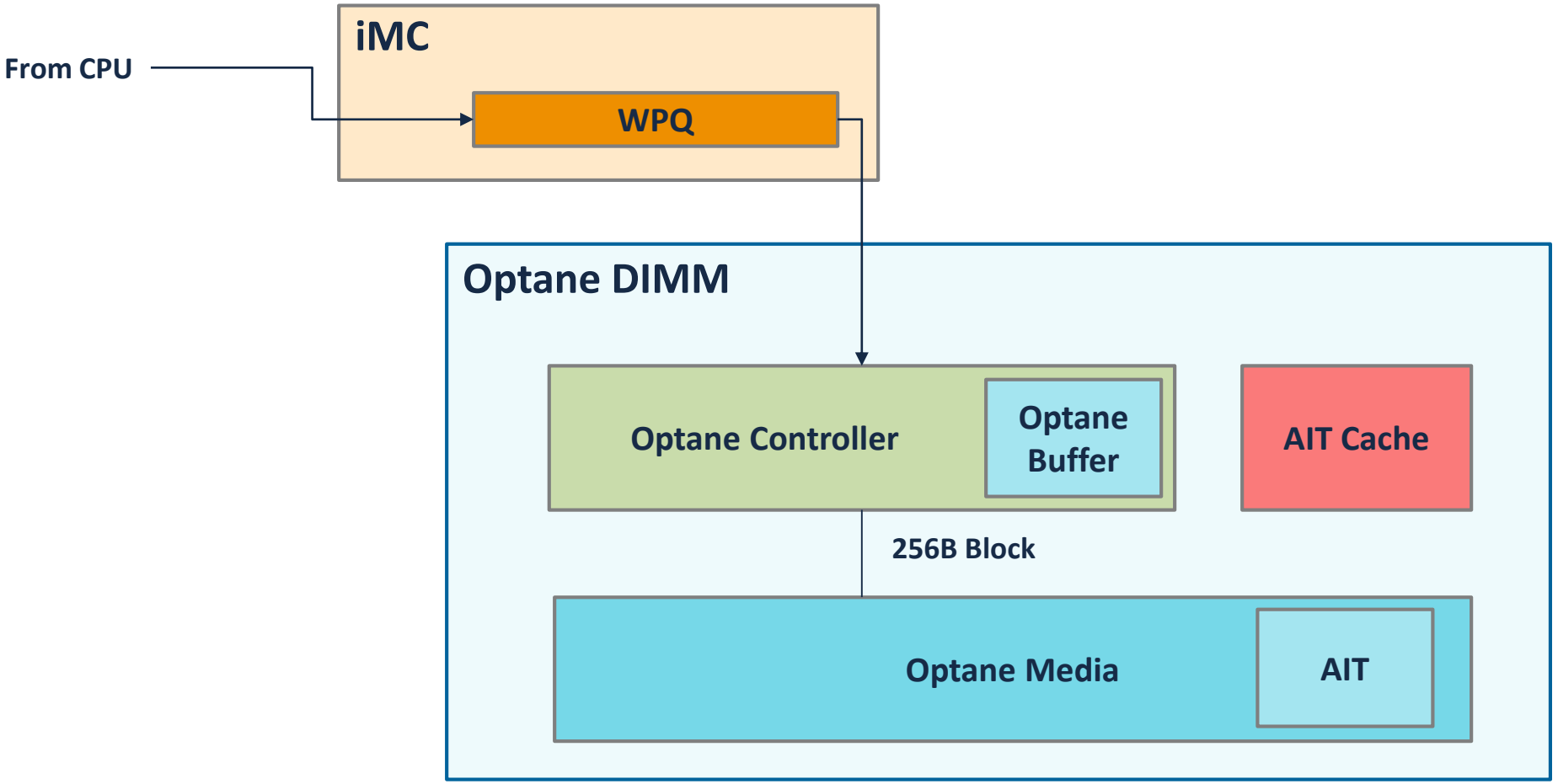

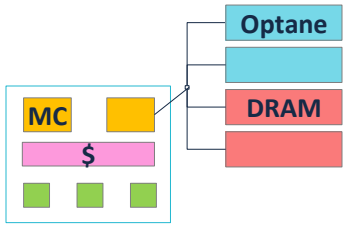

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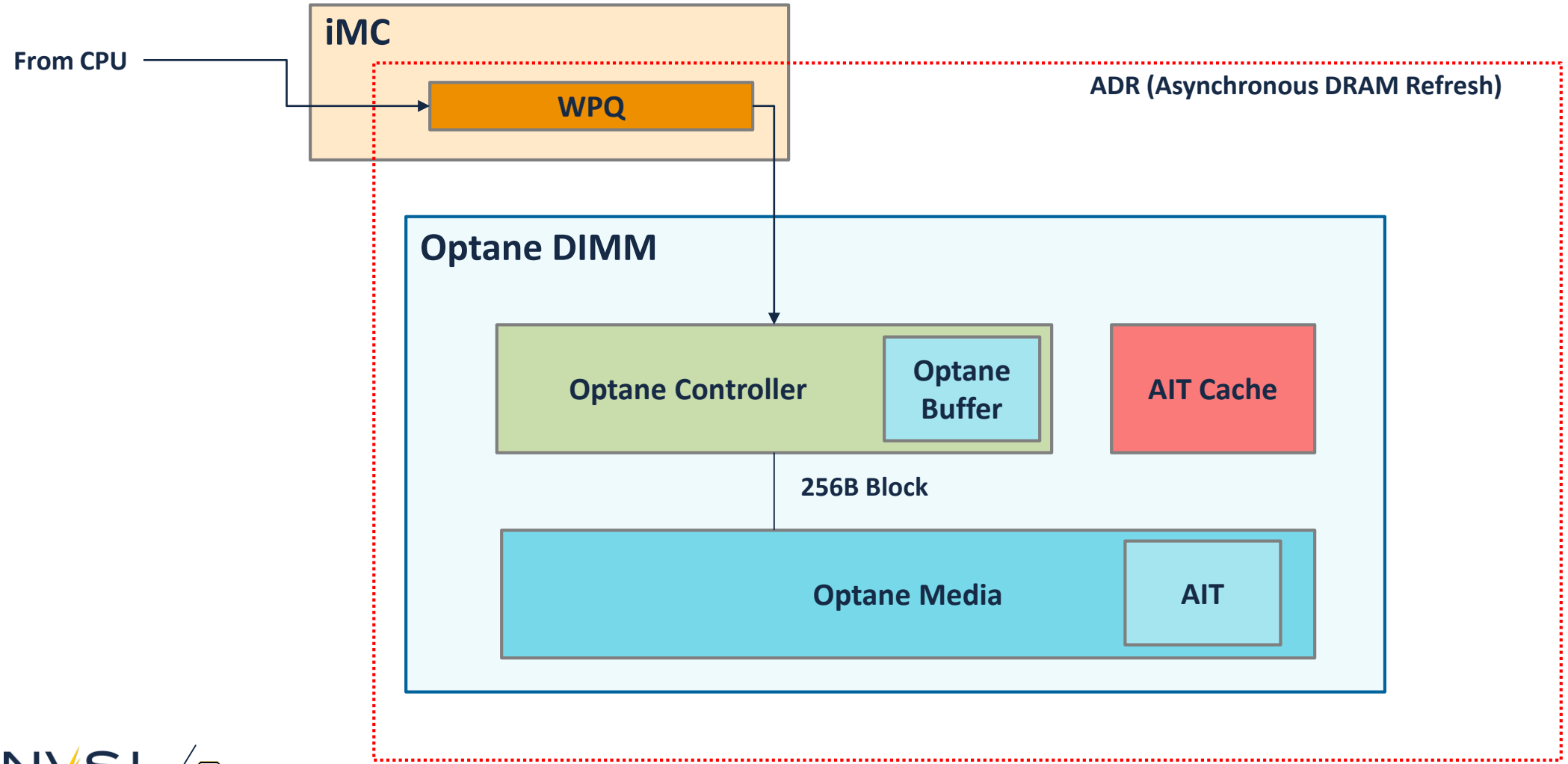
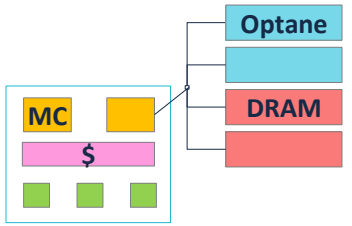
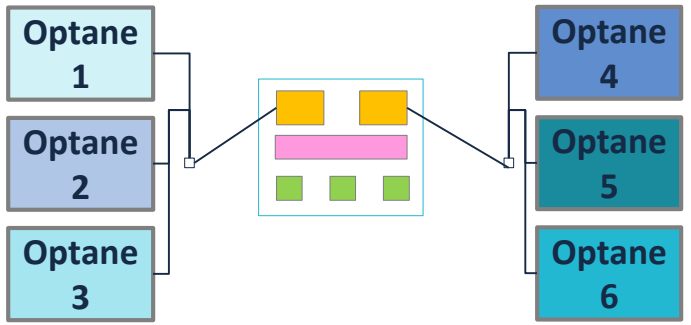




Background: Optane Interleaving

- Optane is interleaved across NVDIMMs at 4KB granularity.



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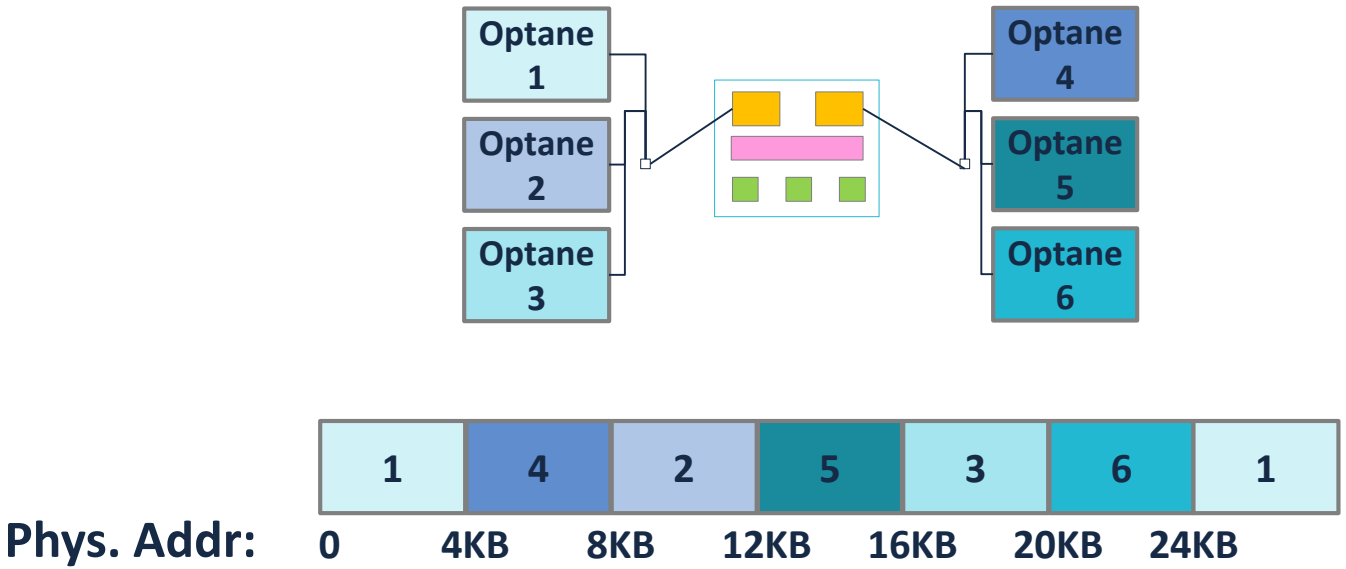


Basics: How does Optane DC perform?

Basics: Our Approach

- Microbenchmark sweeps across state space
  - Access Patterns (random, sequential)
  - Operations (read, ntstore, clflush, etc.)
  - Access/Stride Size
  - Power Budget
  - NUMA Configuration
  - Address Space Interleaving
- Targeted experiments

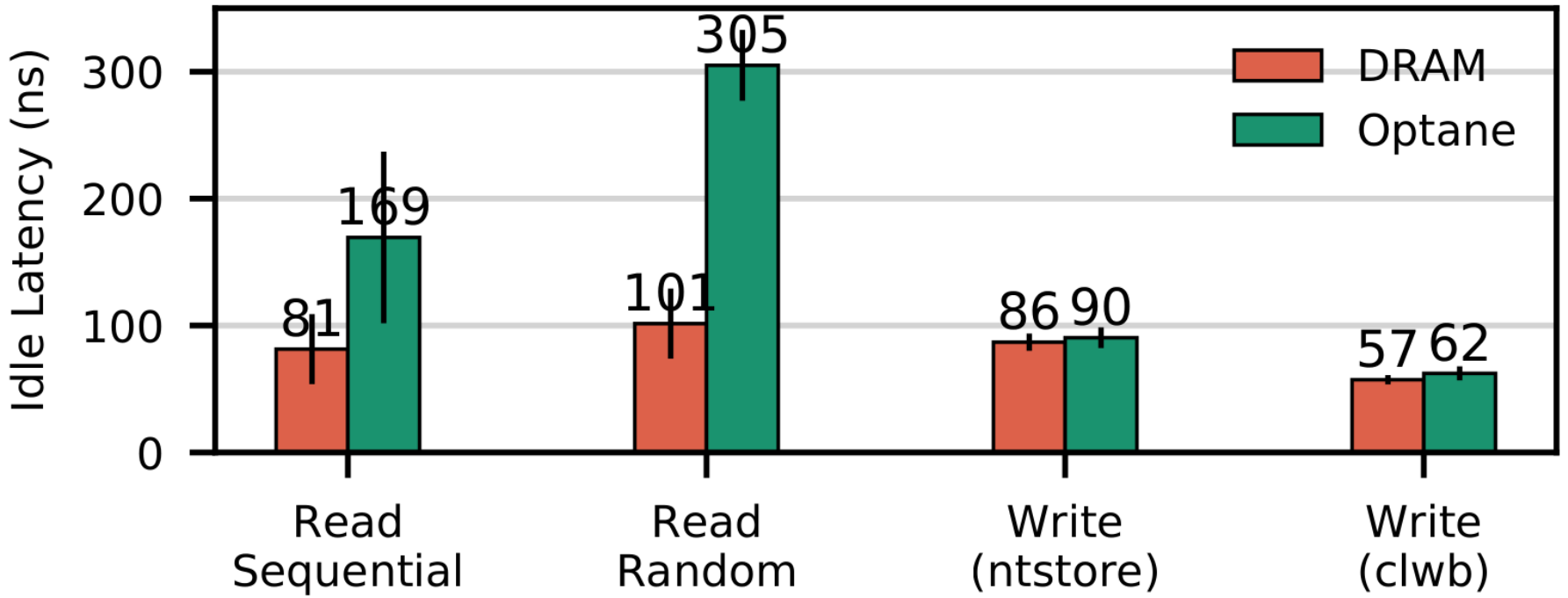
- Total: 10,000 experiments

Test Platform

- CPU
  - Intel Cascade Lake, 24 cores at 2.2 GHz in 2 sockets
  - Hyperthreading off
- DRAM
  - 32 GB Micron DDR4 2666 MHz
  - 384 GB across 2 sockets w/ 6 channels
- Optane
  - 256 GB Intel Optane 2666 MHz QS
  - 3 TB across 2 sockets w/ 6 channels
- OS
  - Fedora 27, 4.13.0

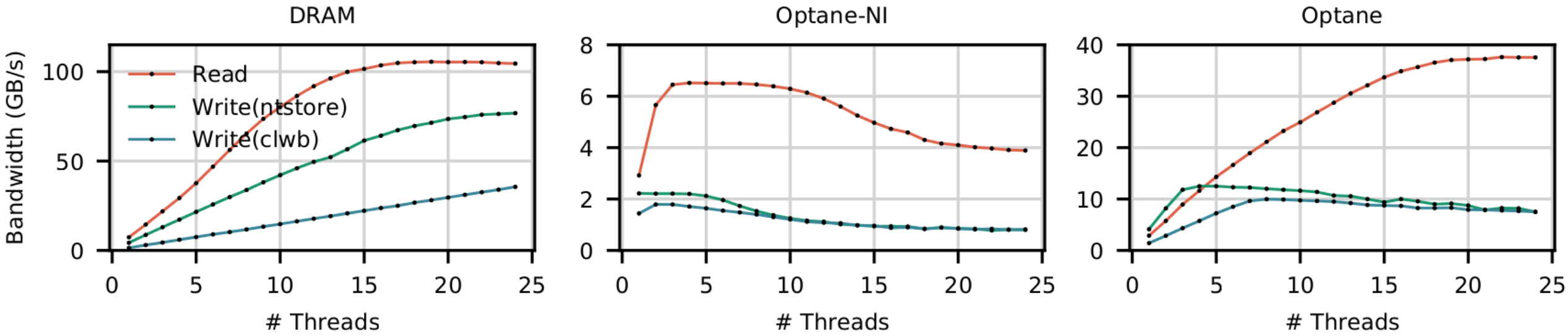
Basics: Latency

- 2x -3x as slow as DRAM
- Write latency masked by ADR



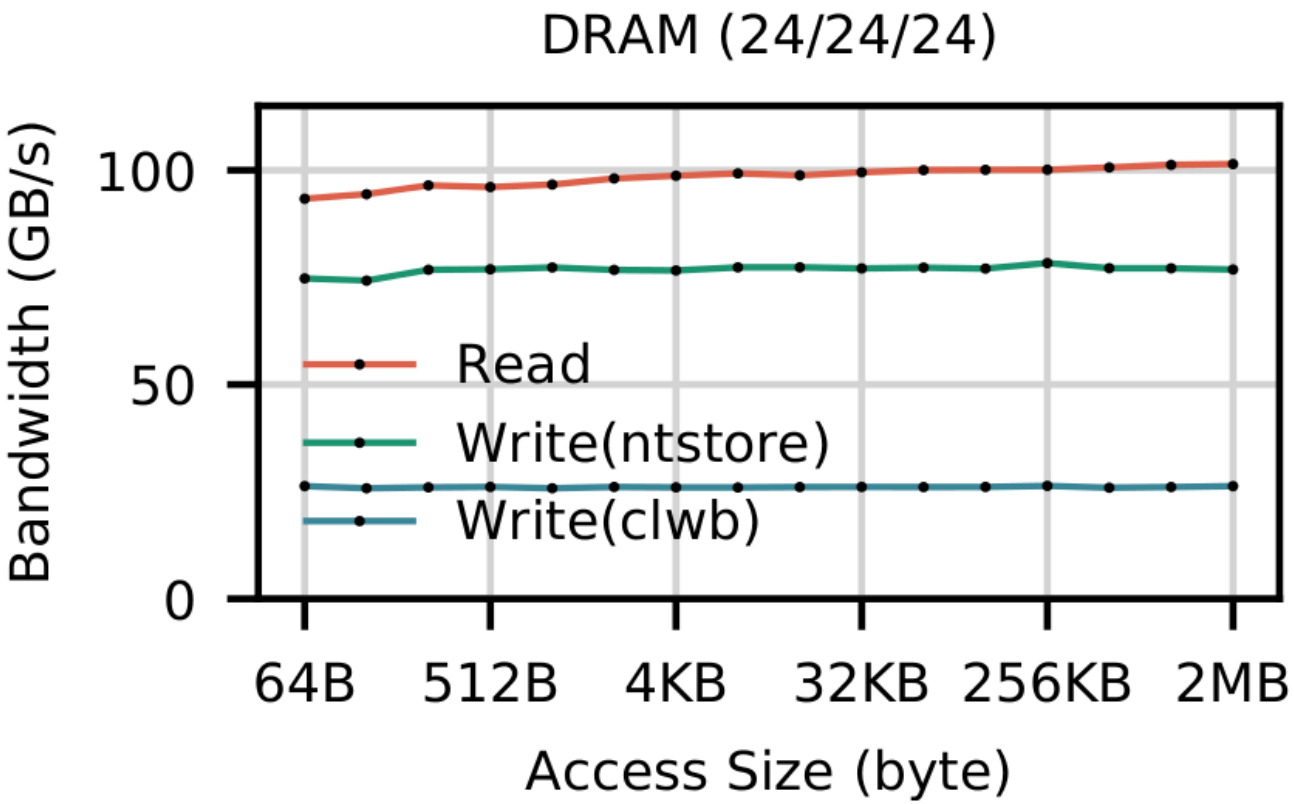
Basics: Bandwidth

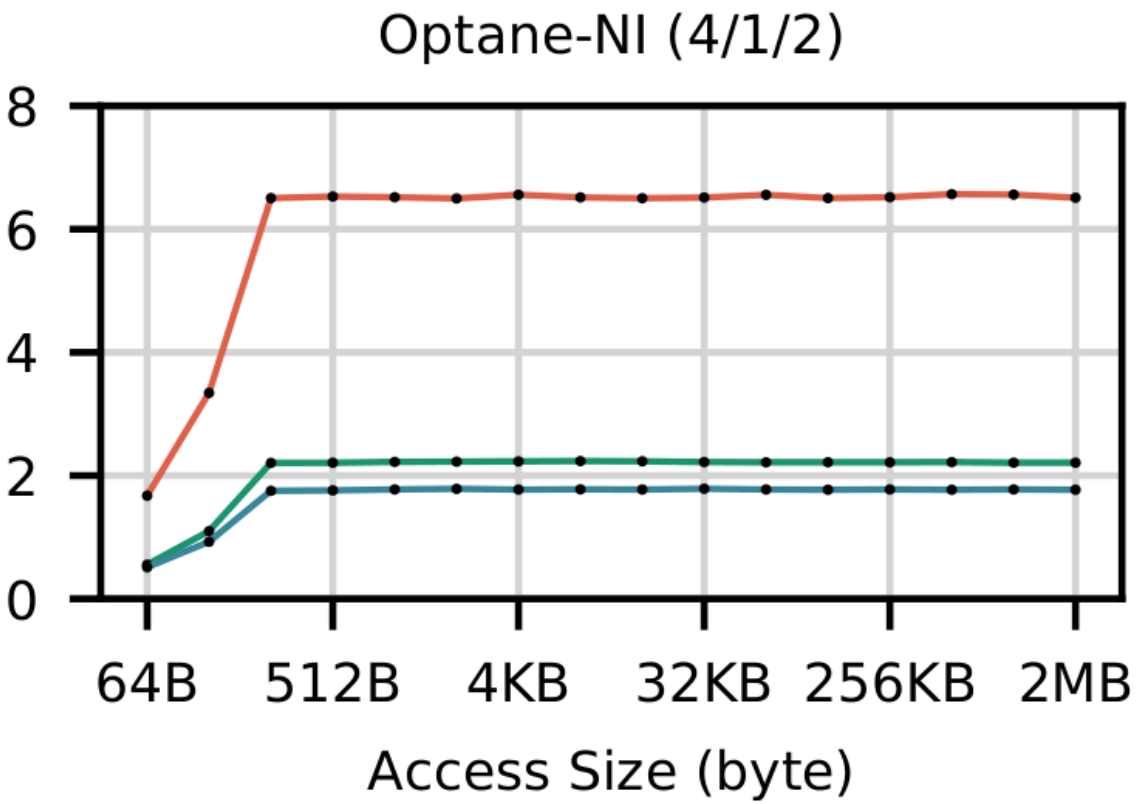
- ~Scalable reads, Non-scalable writes



Basics: Bandwidth

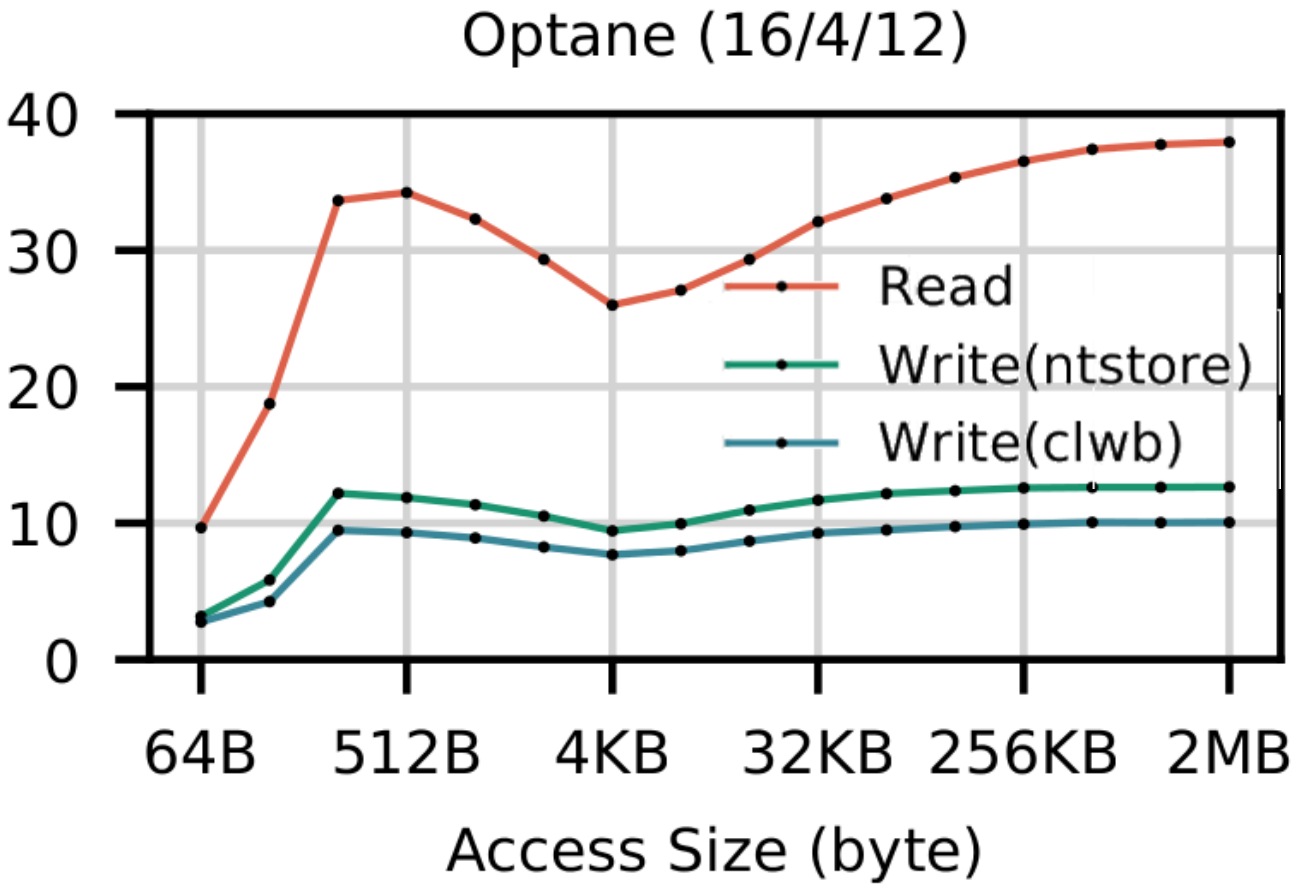
- Access size matters





Basics: Bandwidth

- A mystery!



*answer in ~10 minutes

Lessons: What are Optane Best Practices?

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- Avoid small random accesses
- Use ntstores for large writes
- Limit threads accessing one NVDIMM

Lesson #1: Avoid small random accesses

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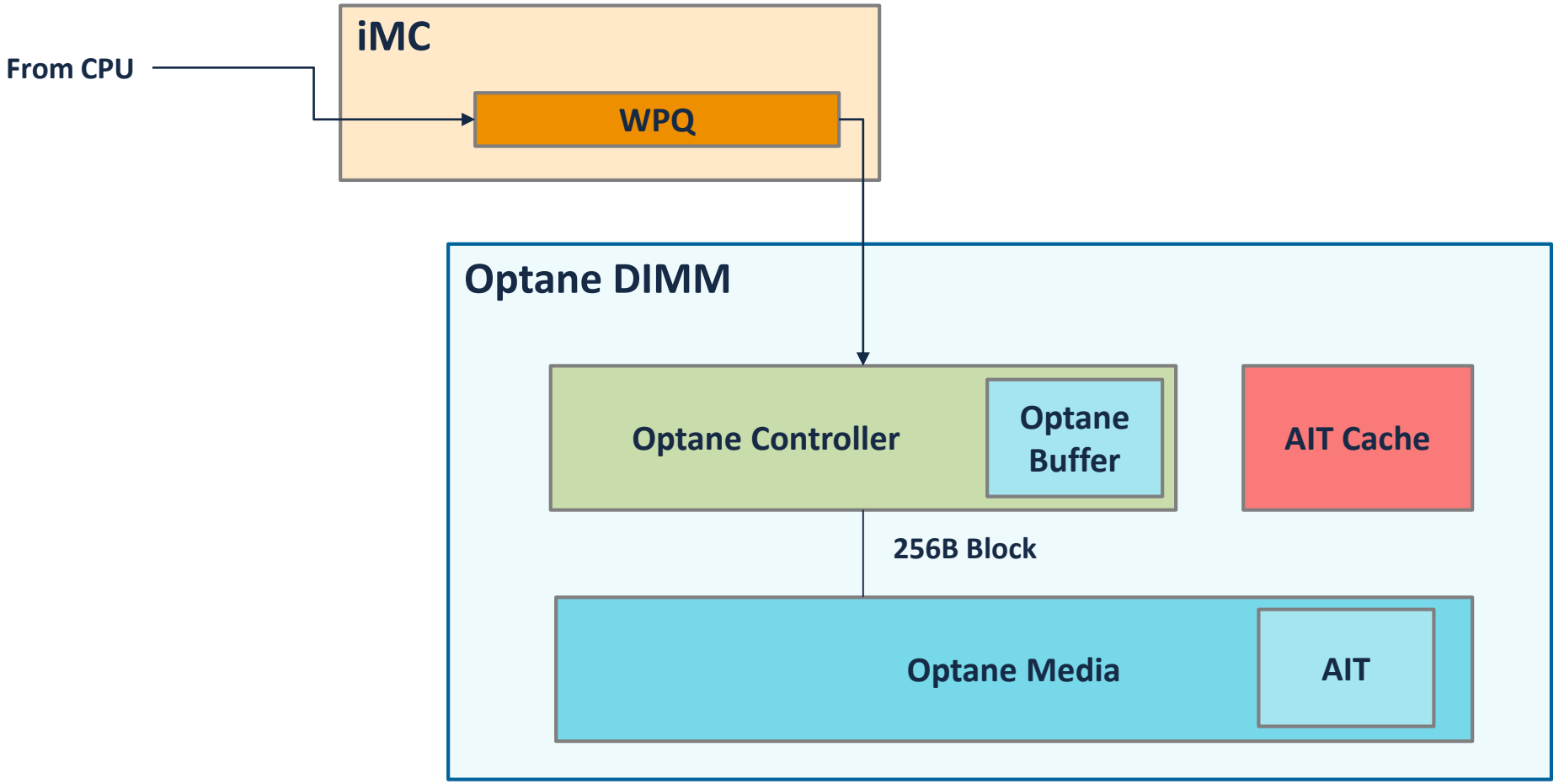

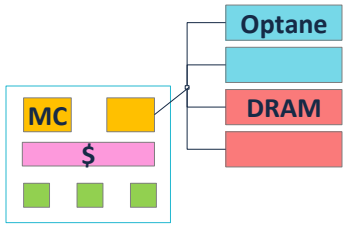

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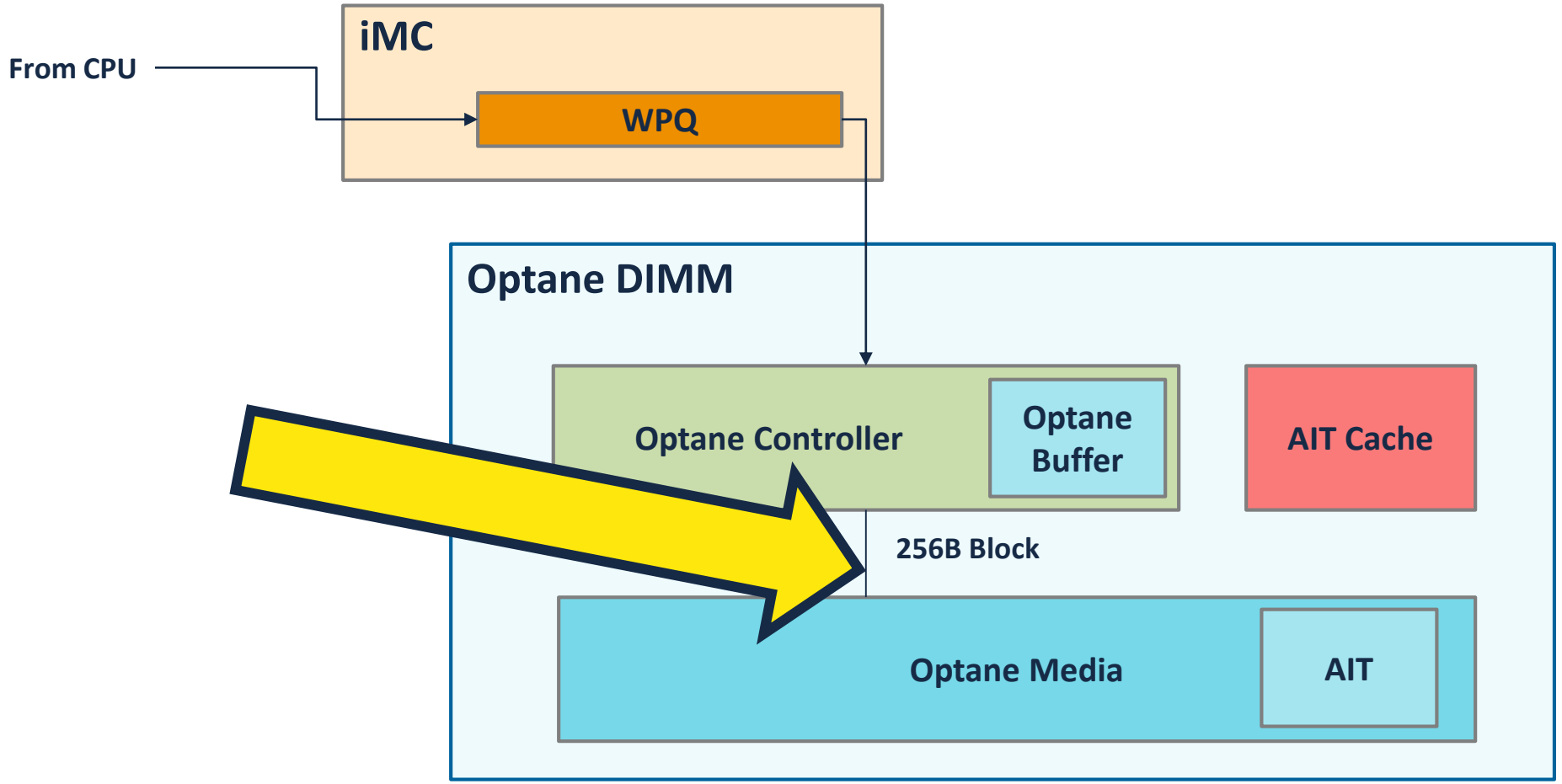

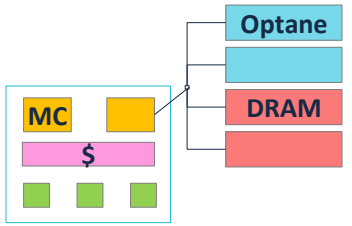

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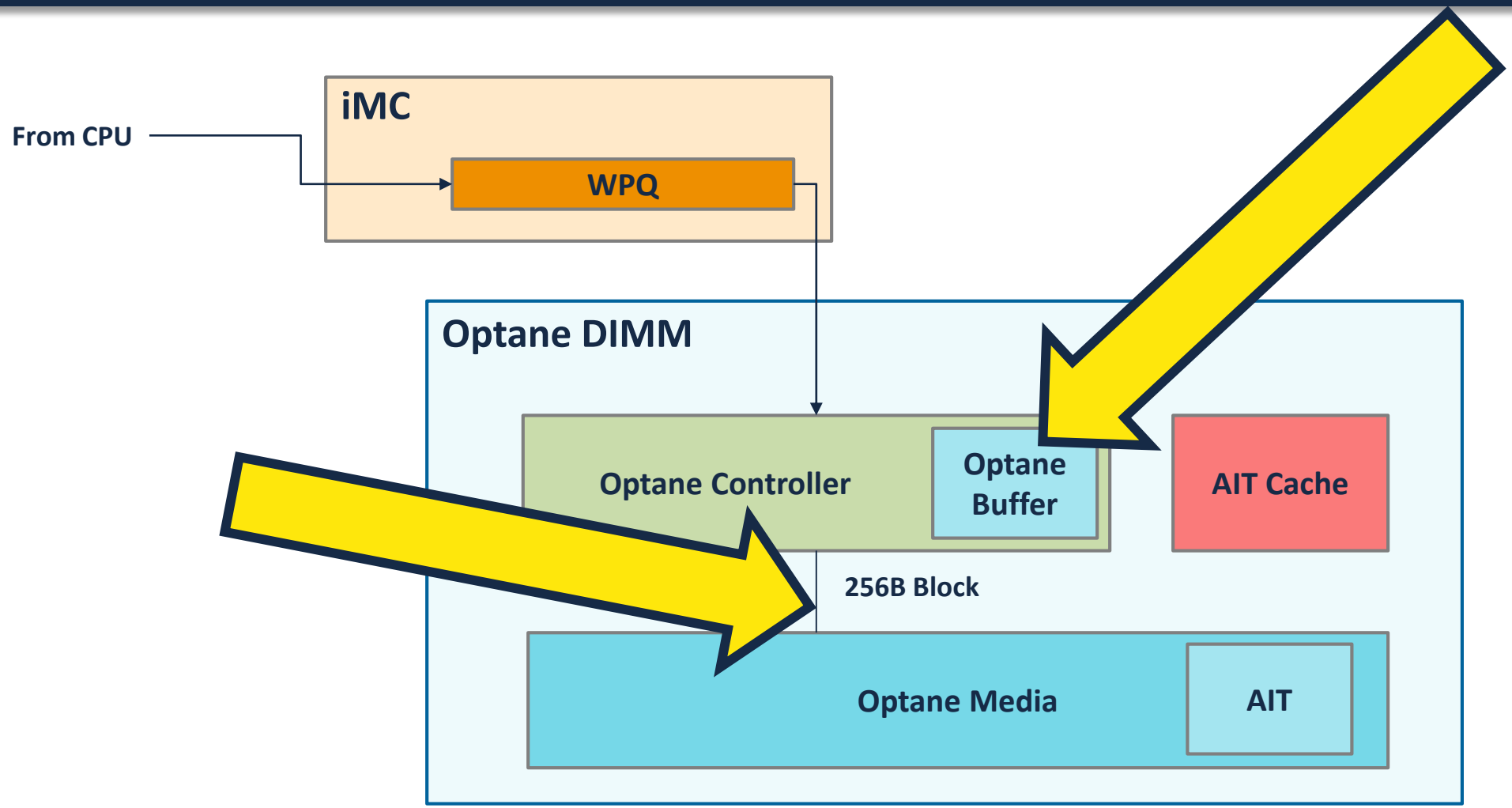
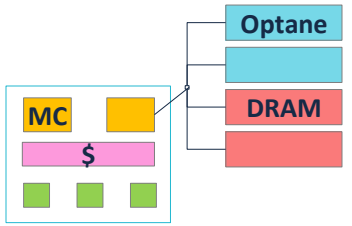
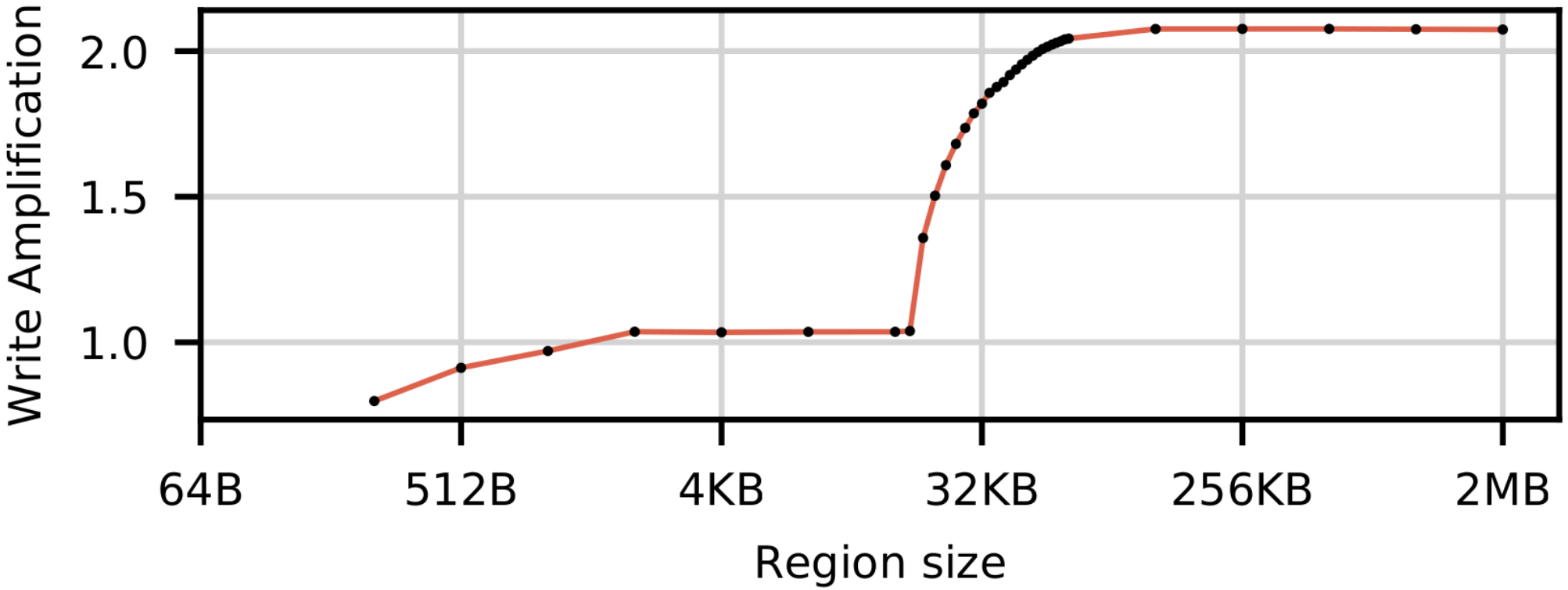




Lessons: Optane Buffer Size

- Write amplification if working set is larger than Optane Buffer



Lesson #1: Avoid small random accesses

- Bad bandwidth with:
  - Small random writes (<256B)
  - Not tiny working set / NVDIMM (>16KB)
- Good bandwidth with
  - Sequential accesses

Lesson #2: Use ntstores for large writes

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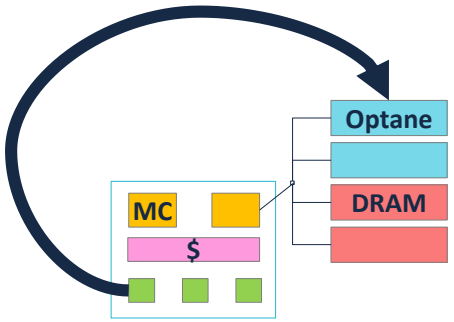
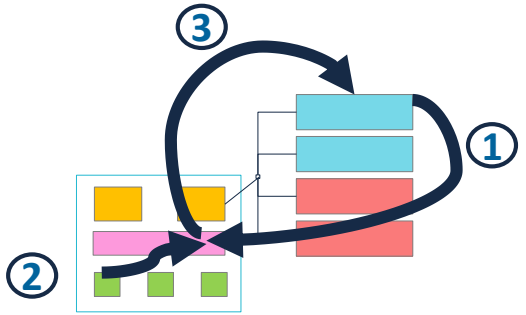
- Non-temporal stores bypass the cache

Lessons: Store instructions

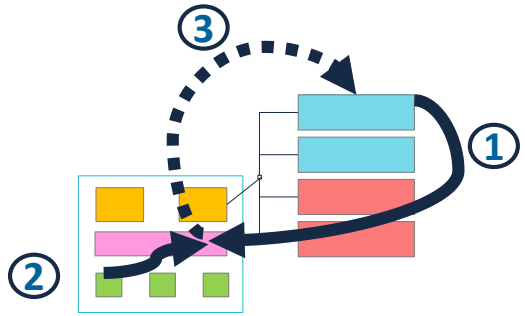
ntstore



store + clwb



store



Lessons: Store instructions

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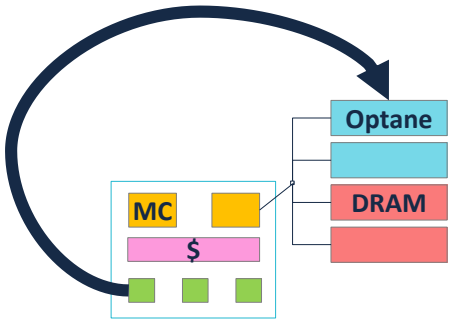
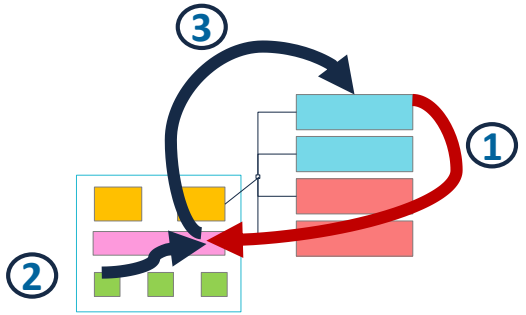
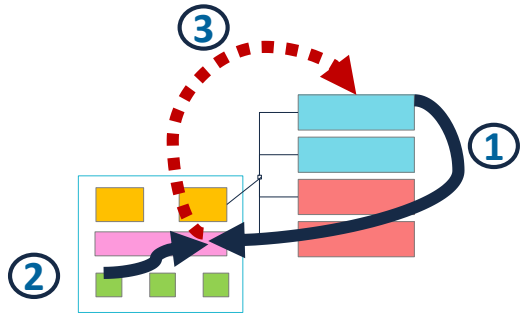


store + clwb



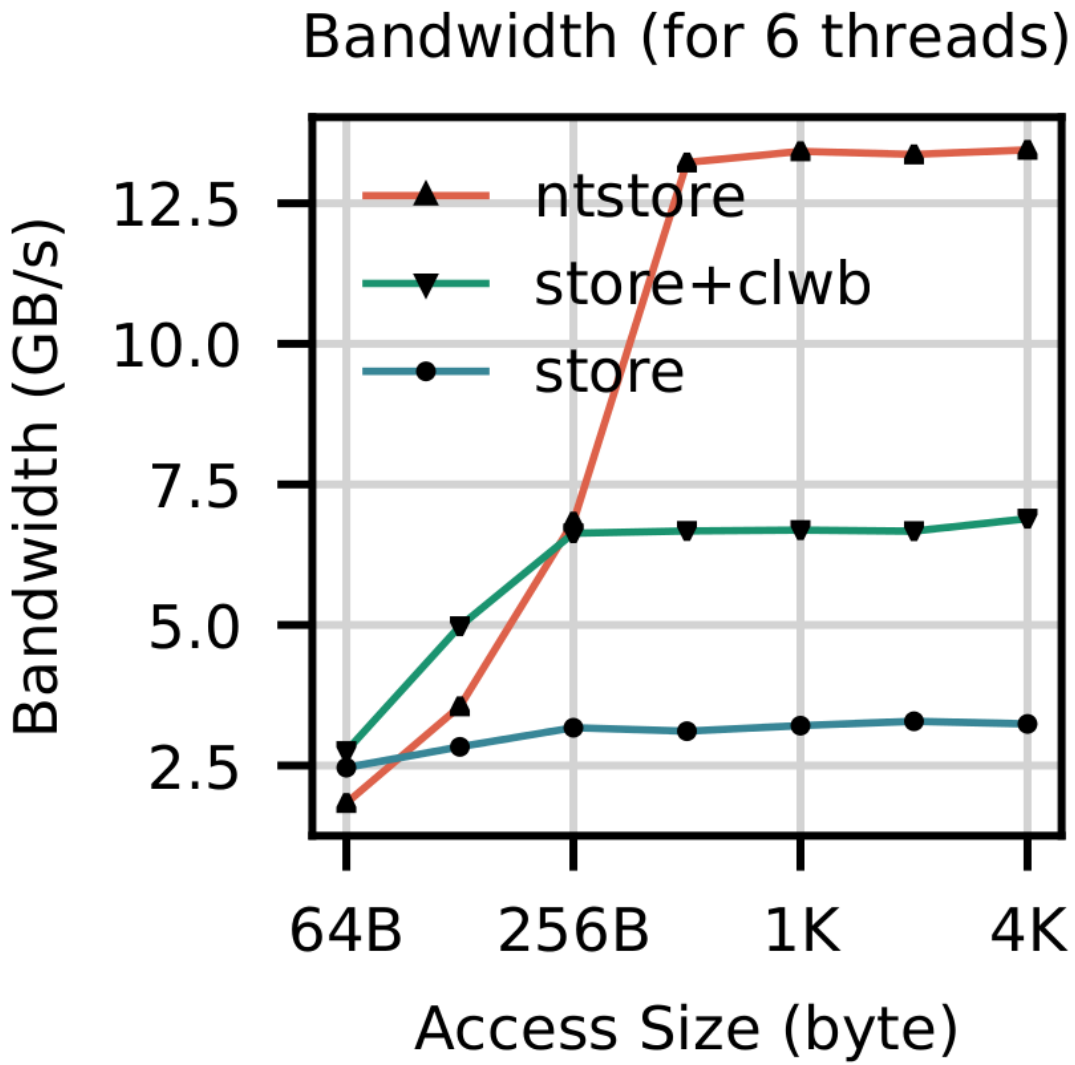
*lost bandwidth

store



*lost locality

Lesson #2: Use ntstores for large writes



Lesson #2: Use ntstores for large writes

- Non-temporal stores bypass the cache
  - Avoid cache-line read
  - Maintain locality

Lesson #3: Limit threads accessing one NVDIMM

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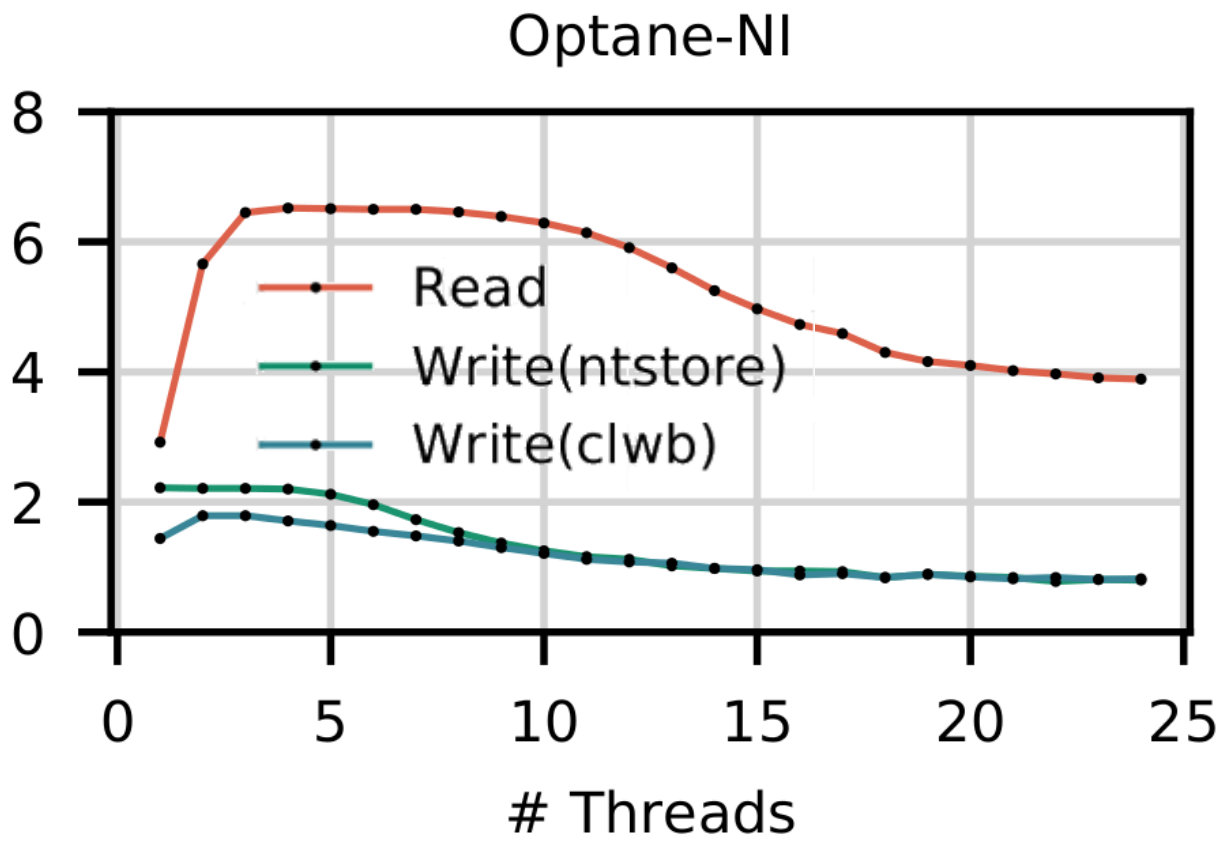
- Contention at Optane Buffer
- Contention at iMC

Lessons: Contention at Optane Buffer

- More threads = access amplification = lower bandwidth

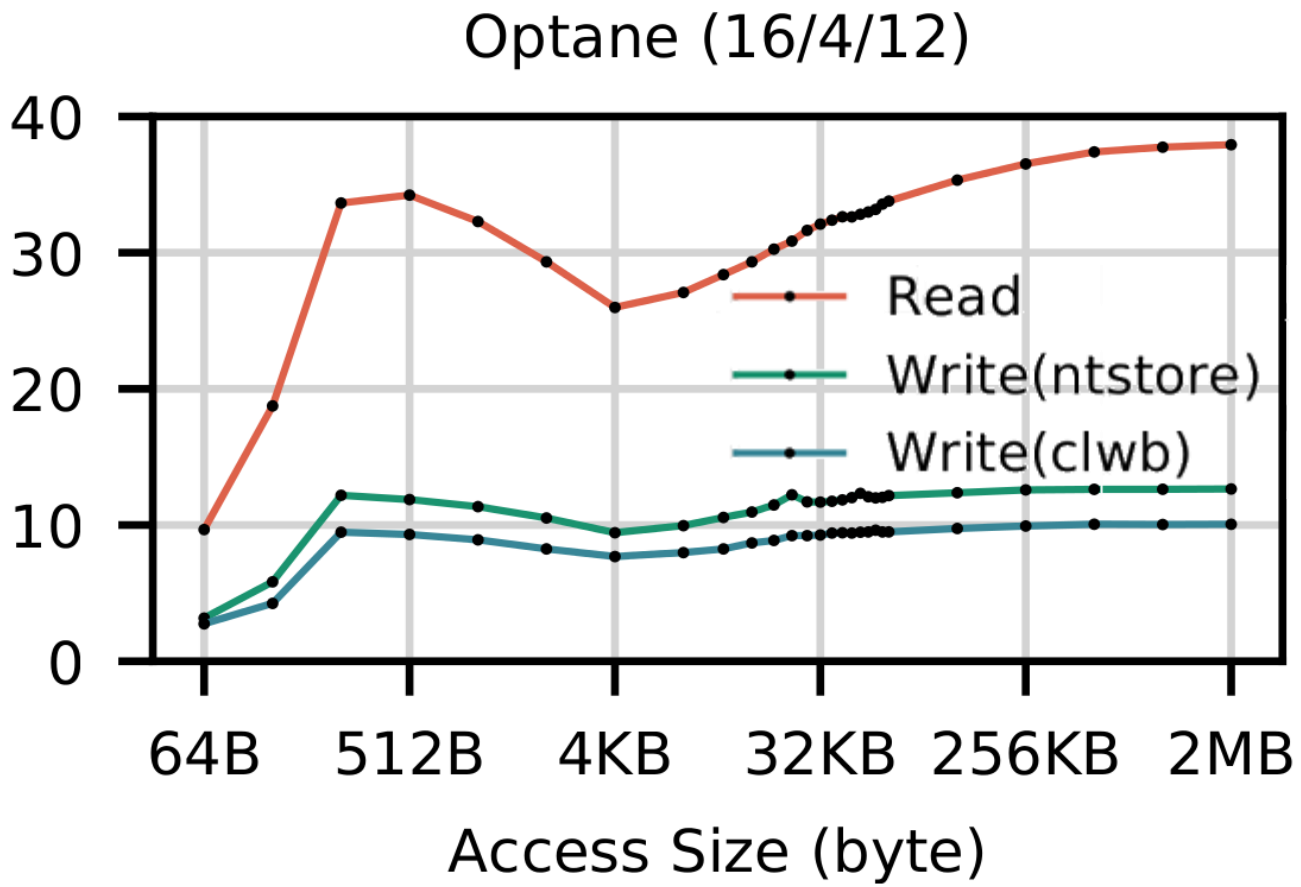


Lessons: Contention at media & iMC

- iMCs queues are small & media is slow
  - Short queues end up clogged

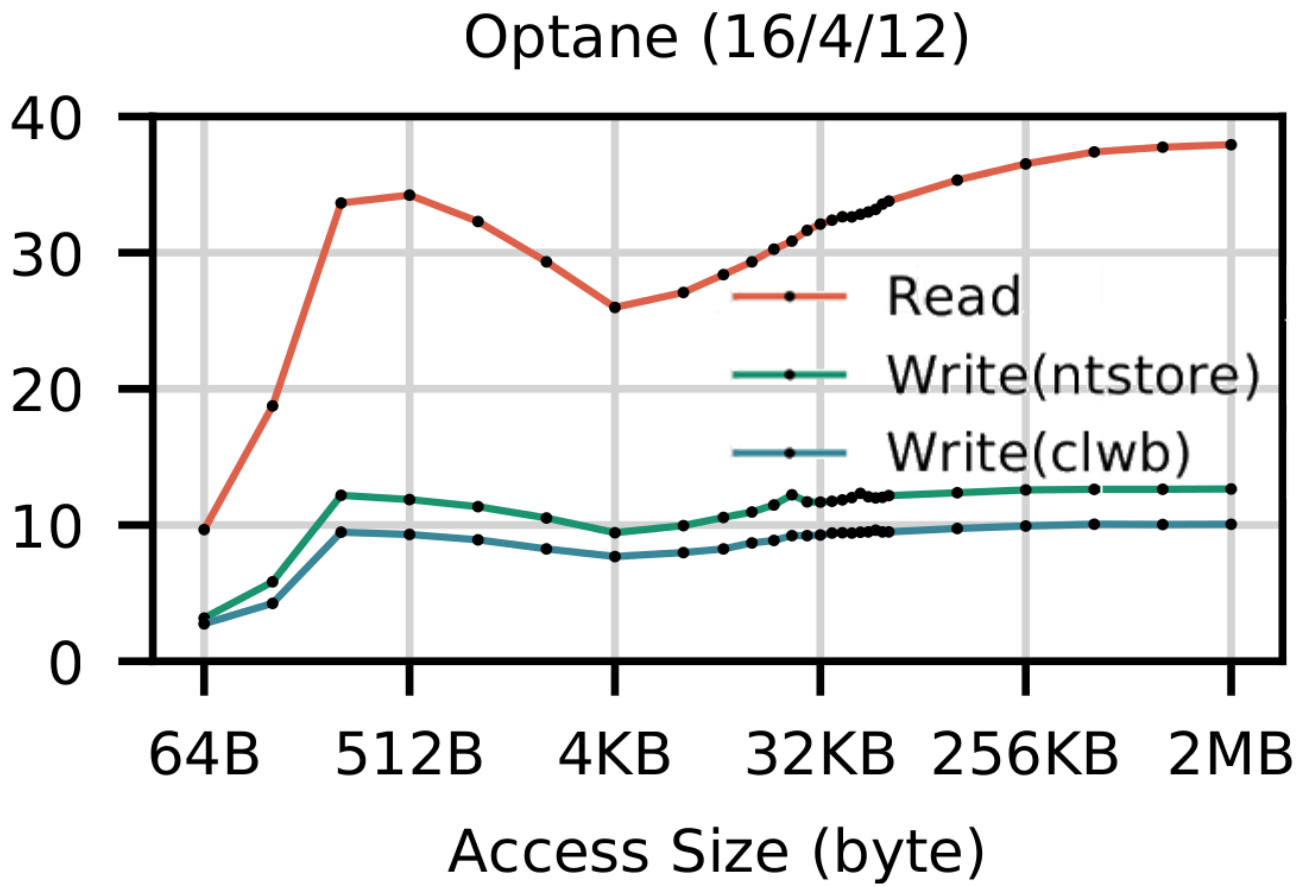
Lessons: Contention at media & iMC

- Contention is largest when random access size = interleave size (4KB)



Lessons: Contention at media & iMC

- Contention is largest when random access size = interleave size (4KB)
- Load is fairest when access size = #DIMMs x interleave size (24KB)



Lesson #3: Limit threads accessing one NVDIMM

- Contention at Optane Buffer
  - Increase access amplification
- Contention at media/iMC
  - Lose bandwidth through uneven NVDIMM load
  - Avoid interleave aligned random accesses

Conclusion

Conclusion

- Not Just Slow Dense DRAM
- Slower media
  -> More complex architecture
    -> Second-order performance anomalies
      -> Fundamentally different

- Max performance is tricky
  - Avoid small random accesses
  - Use ntstores for large writes
  - Limit threads accessing one NVDIMM